

Answer 1: Given $a_0 = 1$, $a_{n+1} = 2a_n + 1$

$$(a) \quad a_0 = 1, \quad a_1 = 3, \quad a_2 = 7,$$

$$a_3 = 15, \quad a_4 = 31$$

$$b) \quad a_n = 2^{n+1} - 1 \quad (\text{Guess})$$

Sanity checks: $a_0 = 1$, $a_1 = 3$, $a_2 = 7 \dots$

(c) Let us apply induction:

Assume $a_n = 2^{n+1} - 1$ is true for n .

① Base case: $a_0 = 2^1 - 1 = 1$ (By definition satisfied)

$$② \quad a_{n+1} = 2a_n + 1$$

$$= 2(2^{n+1} - 1) + 1$$

$$= 2^{n+2} - 1 \quad \text{This implies formula}$$

- - - is true for $n+1$]

$\Rightarrow$ From induction, we have that

$a_n = 2^{n+1} - 1$ is true for all natural numbers.

Answer 2: Existence:

Let $b \geq 2$ be a fixed natural number.

Then from division algorithm, for each $n \in \mathbb{N}$

there exist natural numbers q_0 and a_0 with

$0 \leq a_0 < b$ such that

$$n = q_0 b + a_0$$

If $q_0 = 0$, we are done and $n = a_0$. If not then

we have $n = q_0 b + a_0$. Since $q_0 \neq 0$, we can

apply division algorithm again to get

$$q_0 = q_1 b + a_1 \quad [0 \leq a_1 < b]$$

If $a_1 = 0$, then we are done and

$$n = a_1 b + a_0.$$

Else we have

$$n = a_1 b^2 + a_1 b + a_0$$

This process will end somewhere as naturals are bounded from below by zero.

Let the process ends in $(k+1)$ steps then

$$n = a_0 + a_1 b + \dots + a_k b^k, \text{ here } a_k \neq 0. \\ \text{and } 0 \leq a_i < b \text{ } \forall i.$$

Uniqueness:

$$\text{Suppose } n = a_0 + a_1 b + \dots + a_k b^k$$

$$\text{and } n = c_0 + c_1 b + \dots + c_l b^l$$

$$\text{where } 0 \leq a_i, c_i < b.$$

Assume $k \leq l$ (without loss of generality)

$$\text{Then } 0 = (a_0 - c_0) + \dots + (a_k - c_k) b^k - (c_{k+1} b^{k+1} + \dots + c_l b^l)$$

This expresses 0 as a nontrivial base b expansion.

It is impossible unless $a_i = c_i \quad \forall \quad i \leq k$
and $c_i = 0 \quad \forall \quad i \geq k+1$.

$\Rightarrow$ Uniqueness of base b expansion.

Answer 3: Given $f: \mathbb{Z} \rightarrow \mathbb{Z}$ with $f(x) = x^k$ for $k \geq 2$.

(or) k is a natural number. Then it can either be even or odd.

Case 1 let k is even. Then note that

$$f(1) = 1 = f(-1)$$

$\Rightarrow f$ is not injective.

Case 2 let k is odd. Then $k = 2m+1$ for some

integer $m > 0$ (as $k \geq 2$)

suppose $x, y \in \mathbb{Z}$ and $f(x) = f(y)$. Then

If $x = y$, then function is injective and we are done.

let $x \neq y$. Then either $x > y$ or $x < y$.

Case 1: $x > y$. Put $h = x - y > 0$. Then

$$\begin{aligned} x^k - y^k &= (y+h)^k - y^k \\ &= \sum_{i=1}^k \binom{k}{i} y^{k-i} h^i \end{aligned}$$

(Case 1a): If $y > 0$ then each term

is summation is ≥ 0 and at least

one is > 0 . (eg. $i=k$)

$\Rightarrow x^k - y^k > 0$ Contradicting the assumption.

$$x^k = y^k.$$

$\Rightarrow x$ cannot be greater than y .

(Case 1b): If $y < 0$, but $u = -y$ and $v = -x$

Then $x > y \Rightarrow u > v$.

Since k is odd

$$\begin{aligned} x^k - y^k &= (-v)^k - (-u)^k \\ &= u^k - v^k \end{aligned}$$

Now since $u > v \Rightarrow u^k - v^k > 0$

Again contradicting $x^k = y^k$.

$\Rightarrow x$ can not be greater than y .

Case 2: $x < y$, Reversing the roles of x, y

In case 1, we again get that x cannot

be smaller than y .

$\Rightarrow x = y$.

Thus $f(x) = f(y) \Rightarrow x = y$ when k is odd.

$\Rightarrow f$ is injective function only when k

is odd.

(b) For any $k \geq 2$, not every integer is a perfect k -th power. (e.g. 2 is not power of any integer for any $k \geq 2$)

$\Rightarrow f$ is not surjective.

Hence f is never bijective.

(c) $k=2$, then $f(x) = x^2$

Inverse image of $N = \{0, 1, 2, \dots\}$

$$\begin{aligned} f^{-1}(N) &= \{x \in \mathbb{Z} : x^2 \in N\} \\ &= \mathbb{Z} \end{aligned}$$

(since square of any integer is a natural number)

Answer 4: Given for each rational x

$$d(x) = \min_{m \in \mathbb{Z}} |x - m|$$

(a) By definition of absolute value $|x - m| \geq 0$

$$\Rightarrow d(x) \geq 0 \quad \forall x \in \mathbb{Q}$$

We know that for a rational x , there exists $N \in \mathbb{Z}$ s.t.

$$N \leq x < N+1$$

$$\text{Then } d(x) = \min \{x - N, N+1 - x\}$$

$$\text{Notice that } (x - N) + (N+1 - x) = 1$$

$$\text{Then } \max_{x \in \mathbb{Q}} d(x) = \frac{1}{2}$$

$$\Rightarrow 0 \leq d(x) \leq \frac{1}{2}.$$

(b) Any rational of the form $x = m + \frac{1}{2}$, where $m \in \mathbb{Z}$

$$\text{gives } d(x) = \frac{1}{2}.$$

(c) If $x \in \mathbb{Z}$ then

$$d(x) = \min_{m \in \mathbb{Z}} |x - m|$$

Since $d(x) \geq 0$, the minimum above

is achieved at $m = x$.

$$\Rightarrow d(x)=0 \quad \text{when } x \in \mathbb{Z}.$$

If $d(x)=0$ for some $x \in \mathbb{Q}$. Then

$$d(x) = \min \{x-N, N+1-x\}$$

But $\min \{x-N, N+1-x\} = 0$ only if either

$$x=N \quad \text{or} \quad x=N+1.$$

That only if x is an integer.

$$\Rightarrow d(x)=0 \quad \text{iff } x \in \mathbb{Z}.$$

Answer 5: Claim: No surjection from $f: \mathbb{N} \rightarrow \mathcal{P}(\mathbb{N})$

Proof: Suppose for contradiction that such a surjection exists. This means that for each $A \subseteq \mathbb{N}$, there exists some $n \in \mathbb{N}$ s.t.

$$f(n) = A.$$

Construct the set

$$S = \{n \in \mathbb{N} : n \notin f(n)\}$$

Note that S is a subset of $\mathbb{N}$ by construction.

That is $S \subseteq \mathbb{N}$. Then there is some $k \in \mathbb{N}$ s.t.

$$f(k) = S. \quad [\text{As } f \text{ is surjective}]$$

Now ask if $k \in S$?

① If $k \in S$. Then by definition of S ,

$$k \notin f(k) \Leftrightarrow k \notin S.$$

[A contradiction]

② If $k \notin S$. Then again $k \notin f(k)$ as $f(k) = S$.

Then by definition of S , $k \in S$.

[A contradiction]

$\Rightarrow$ A surjective $f: \mathbb{N} \rightarrow \mathcal{P}(\mathbb{N})$ cannot exist.

$$\Rightarrow |\mathbb{N}| < |\mathcal{P}(\mathbb{N})|$$